Let's write a LISP lexer together

Coding dojo by @meisterluk





Vocabulary

- An *interpreter* reads source code and applies semantics immediately
- A formal *grammar* specifies the set of admissible source codes for the interpreter
- An interpreter can include a lexical analysis / tokenization (component "*lexer*") and a semantic analysis / parsing (component "*parser*")

Input source code:

W

Input source code:

W

Output:

Hello world!

Input source code:

Admissible source codes:

W

W

Output:

Hello world!

Input source code:

Ρ

Output:

Hello pygraz!

Input source code:

Admissible source codes:

W P

Output:

Input source code:

Admissible source codes:

W P

Output:

Implementation:

```
if src == "W":
    print("Hello world!")
else:
    print("Hello pygraz!")
```

Input source code:

Admissible source codes:

W P

Output:

Implementation:

if src.strip() == "W":
 print("Hello world!")
else:
 print("Hello pygraz!")

Boooring – can we do something non-static?

A non-static formal grammar

Input source code:

put Hello world!

Admissible source codes:

put <some-string>

Output:

Hello world!

A non-static formal grammar

Input source code:

put Hello world!

Admissible source codes:

put <some-string>

Output:

Hello world!

Implementation:

assert(src[0:4] == "put ")
print(src[4:])

But what if we need to compute the output string beforehand?

Input source code:

put Hello sum(4, 5)!

Admissible source codes:

put <some-string>
and expression sum(<args>)

Output:

Implementation:

Hello 9!

Input source code:

put Hello sum(4, 5)!

Admissible source codes:

put <some-string>
and expression sum(<args>)

Output:

Hello 9!

Implementation:

?

put Hello sum(4, 5)!

identifiers arguments operators

put Hello sum(4, 5)!

identifiers arguments operators

Why are identifiers invoked so differently?

Where do I need commas between arguments?

What happens if I use "sum" as argument?

The simplest nested formal grammar

(put Hello (sum 4 5) !)

identifiers arguments operators

The simplest nested formal grammar

(put Hello (sum 4 5) !)

identifiers arguments operators

parenthesized prefix notation

The simplest nested formal grammar

(put Hello (sum 4 5) !)

identifiers arguments operators

parenthesized prefix notation = LISP?

LISP

Timeline [edit]

					Timeline of Lisp dialects				
1958	1960	1965	1970	1975	1980	1985	1990	1995	
LISP 1, 1 2(abandone									
		Maclisp							
			Interlisp						
			MDL						
				Lisp Mad	hine Lisp				
				Scheme				R5RS	
				NIL					
					ZIL (Zori				
					Implemen				
					Language				
					Franz Lis	-			
					Common	Lisp	ANSI sta	ndard	
					Le Lisp				
					MIT Sche	eme			
					XLISP				

1985	1990	1995	2000	2005	2010	2015	2020
т							
Chez Sch	eme						
Emacs Li	sp						
AutoLISP	•						
PicoLisp							
Gambit							
	EuLisp						
	ISLISP						
	OpenLisp)					
	PLT Sche	me			Racket		
	newLISP						
	GNU Guil	е					
			Visual LI	SP			
				Clojure			
				Arc			
				LFE			
					Ну		
							Chialisp

```
#lang racket/base
```

```
(provide ~> ~>> and~> and~>> _
```

```
\begin{split} & |\operatorname{ambda} > |\operatorname{ambda} >> |\operatorname{ambda} >> |\operatorname{ambda} ->> * \\ & |\operatorname{ambda} - \operatorname{and} >> |\operatorname{ambda} - \operatorname{and} ->> |\operatorname{ambda} - \operatorname{and} ->> * \\ & (\operatorname{rename-out} [|\operatorname{ambda} -> \lambda ->] [|\operatorname{ambda} ->> \lambda ->>] \\ & [|\operatorname{ambda} -> * \lambda ->*] [|\operatorname{ambda} ->> * \lambda ->>*] \\ & [|\operatorname{ambda} - \operatorname{and} ->> * ||\operatorname{ambda} - \operatorname{and} ->> || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda} - \operatorname{and} ->> * || || \\ & [|\operatorname{ambda}
```

(defn- alias-help

```
=> (print "Hy!")
Hy!
=> (defn salutationsnm [name] (print (+ "Hy " name "!")))
=> (salutationsnm "YourName")
Hy YourName!
```

```
(define (make-account)
 (let ((balance 0))
   (define (get-balance)
      balance)
   (define (deposit amount)
      (set! balance (+ balance amount))
      balance)
   (define (withdraw amount)
      (deposit (- amount)))
```

```
(lambda args
  (apply
    (case (car args)
        ((get-balance) get-balance)
        ((deposit) deposit)
        ((withdraw) withdraw)
        (else (error "Invalid method!")))
        (cdr args)))))
```

In my opinion

- *LISP* is a programming language with syntax (defining a formal grammar) and semantics
- The LISP family/dialects is a set of programming languages following the style of LISP 1.0 or LISP 1.5
- The syntax of LISP is called *S*-expressions.
- S-expressions is a form of *parenthesized prefix notation*

S-expressions

GNU Guile 3.0.7 Copyright (C) 1995-2021 Free Software Foundation, Inc. Guile comes with ABSOLUTELY NO WARRANTY; for details type `,show w'. This program is free software, and you are welcome to redistribute it under certain conditions; type `,show c' for details. Enter `,help' for help. scheme@(guile-user)> (cons 1 3) \$1 = (1 . 3) scheme@(guile-user)> (cons 1 (cons 3 '())) \$2 = (1 3) scheme@(guile-user)>

Characteristics [edit]

In the usual parenthesized syntax of Lisp, an S-expression is classically defined^[1] as

- 1. an atom of the form x, or
- 2. an expression of the form $(x \cdot y)$ where x and y are S-expressions.

This definition reflects LISP's representation of a list as a series of "cells", each one an ordered pair. In plain lists, *y* points to the next cell (if any), thus forming a list. The recursive clause of the

Coding Dojo

Task: Let us read a file written in parenthesized prefix notation.

Coding Dojo

Task:

Let us read a file written in parenthesized prefix notation.

Funfacts:

- The standard library code module provides a REPL to parse python code: https://bernsteinbear.com/blog/simple-python-repl/
- Peter Norvig documented our task in a blog post: https://norvig.com/lispy.html
 "The beauty of Scheme is that the full language only needs 5 keywords and 8 syntactic forms. In comparison, Python has 33 keywords and 110 syntactic forms, and Java has 50 keywords and 133 syntactic forms."

Approach

- 1) Identify individual characters of the formal grammar. Give them names.
- 2) Define an INIT state. Which characters are admissible?
- 3) Reiterate to identify the lexing state diagram.
- 4) Yield tokens as you read character by character
- 5) A parser fetches these tokens to put them into a nested structure
- 6) The nested structure is interpreted.

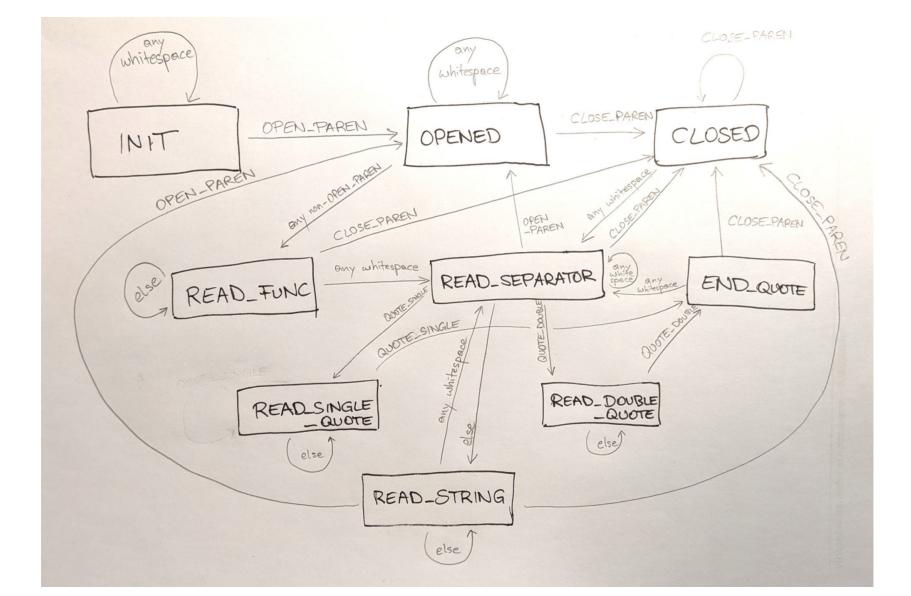
Approach

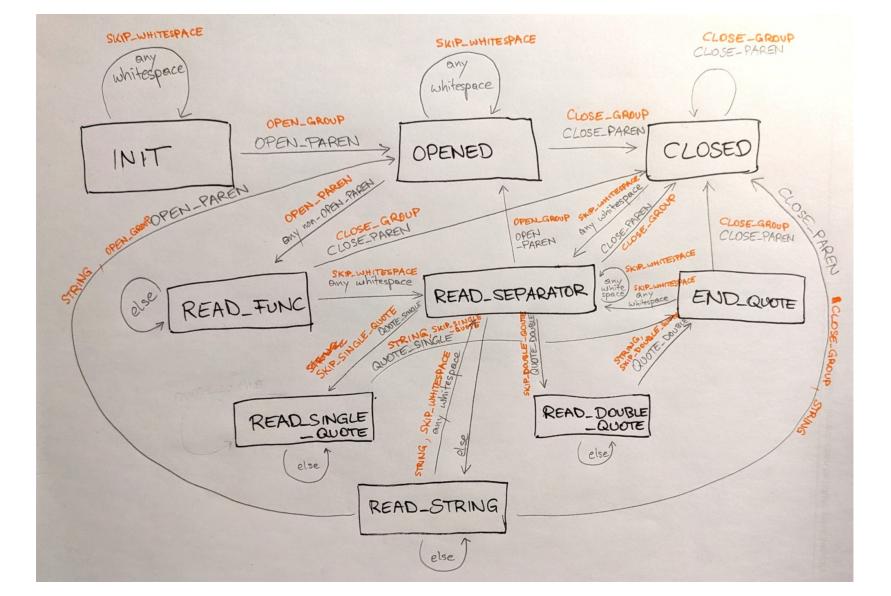
- 1) Identify individual characters of the formal grammar. Give them names.
- 2) Define an INIT state. Which characters are admissible?
- 3) Reiterate to identify the lexing state diagram.
- 4) Yield tokens as you read character by character
- 5) A parser fetches these tokens to put them into a nested structure

6) The nested structure is interpreted.

parser

lexer





```
# raw characters of the syntax, I want to match
OPEN_PAREN = '('
CLOSE_PAREN = ')'
QUOTE_SINGLE = "'"
QUOTE_DOUBLE = '"'
```

class LexingState(enum.Enum):

```
INIT = 1
OPENED = 2
CLOSED = 3
READ_FUNC = 4
READ_SEPARATOR = 5
READ_STRING = 6
READ_SINGLE_QUOTE = 7
READ_DOUBLE_QUOTE = 8
END_QUOTE = 9
```

def main(spec_file):

"""Read the specification file and represent it in an arbitrary way so the user can verify that the file is interpreted appropriately. """

lex = LispLexer(spec_file)

```
with open(spec_file) as fd:
    content = fd.read()
    par = LispParser(content)
```

```
for token in lex.feed(content):
    tok, start, end = token
    print(tok, repr(content[start:end]))
```

```
par.consume_token(token)
```

```
spec_tree = par.finalize()
```

```
print()
print(repr(spec_tree))
```

class LispLexer:

```
def __init__(self, source):
    self.state = LexingState.INIT
    self.scalar_id = 0
    self.column_id = 0
    self.line_id = 0
    self.start_string = None
    self.source = source
```

class LexedToken(enum.Enum): OPEN_GROUP = 1 # start a new group STRING = 2 # gives a new string argument CLOSE_GROUP = 3 # terminates a group SKIP_WHITESPACE = 4 # whitespace not required to understand semantics SKIP_SINGLE_QUOTE = 5 # single quote not required to understand semantics SKIP_DOUBLE_QUOTE = 6 # double quote not required to understand semantics

Let's go!